# 1 Math

**Defn.** Even and Odd Integers

n is even  $\Leftrightarrow \exists$  an integer k s.t. n=2kn is odd  $\Leftrightarrow \exists$  an integer k s.t. n=2k+1

**Defn.** Divisibility

n and d are integers and  $d \neq 0$ 

 $d|n \Leftrightarrow \exists k \in \mathbb{Z} \text{ s.t. } n = dk$ 

Theorem 4.2.1. Every Integer is a rational number

Theorem 4.2.2. The sum of any two rational numbers is rational

**Theorem 4.3.1.** For all  $a, b \in \mathbb{Z}^+$ , if a|b, then  $a \leq b$ 

**Theorem 4.3.2.** Only divisors of 1 are 1 and -1

**Theorem 4.3.3.**  $\forall a, b, c \in \mathbb{Z} \text{ if } a|b, b|c, a|c$ 

Theorem 4.6.1. There is no greatest integer

**Proposition.** 4.6.4 For all integers n, if  $n^2$  is even, then n is even.

**Defn.** Rational r is rational  $\Leftrightarrow \exists a, b \in \mathbb{Z} \text{ s.t. } r = \frac{a}{b} \text{ and } b \neq 0$ 

**Defn.** Fraction in lowest term: fraction  $\frac{a}{b}$  is lowest term if largest  $\mathbb{Z}$  that divies both a and b is 1

**Theorem 4.7.1.**  $\sqrt{2}$  is irrational

# 2 Logic of Combound Statements

**Theorem 3.2.1.** Negation of universal stmt  $\sim (\forall x \in D, P(x)) \equiv \exists x \in D \text{ s.t. } \sim P(x)$ 

**Theorem 3.2.1.** Negation of existential stmt  $\sim (\exists x \in D \text{ s.t. } P(x)) \equiv \forall x \in D, \sim P(x)$ 

**Defn.** Contrapositive of  $p \Rightarrow q \equiv \sim q \Rightarrow p$ 

**Defn.** Converse of  $p \Rightarrow q$  is  $q \Rightarrow p$ 

**Defn.** Inverse of  $p \Rightarrow q$  is  $\sim p \Rightarrow \sim q$ 

**Defn.** Only if: p only if q means  $\sim q \Rightarrow \sim p \equiv p \Rightarrow q$ 

**Defn.** Biconditional:  $p \Leftrightarrow q \equiv (p \Rightarrow q) \land (q \Rightarrow p)$ 

**Defn.** r is sufficient condition for s means if r then  $s, r \Rightarrow s$ 

**Defn.** r is necessary condition for s means if  $\sim r$  then  $\sim s, s \Rightarrow r$ 

**Defn.** Proof by Contradiction

If you can show that the supposition that st<br/>tatement p is false leads to a contradiction, then you can conclude that p is true

# 3 Methods of Proof

Statement	Proof Approach		
$\forall x \in D \ P(X)$	Direct: Pick arbitrary x, prove P is true for that x.		
	Contradiction: Suppose not, i.e. $\exists x (\sim p)$ Hence supposition $\sim p$ is false (P3)		
$\exists x \in D \ P(X)$	Direct: Find x where P is true.		
	Contradiction: Suppose not, i.e. $\forall x (\sim p)$ Hence supposition $\sim p$ is false (P3)		
$P \Rightarrow Q$	Direct: Assume P is true, prove Q		
	Contradiction: Assume P is true and Q is false, then derive contradiction		
	Contrapositive: Assume $\sim Q$ , then prove $\sim P$		
$P \Leftrightarrow Q$	Prove both $P \Rightarrow Q$ and $Q \Rightarrow P$		
xRy. Prove R is equivalence	Prove Reflexive, Symmetric and Transitive		
Reflexive			
Symmetric			
Antisymmetric			
Transitive			

**Defn.** Proof by Contraposition

- 1. Statement to be proved  $\forall x \in D \ (P(x) \Rightarrow Q(x))$
- 2. Contrapositive Form:  $\forall x \in D \ (\sim Q(x) \Rightarrow \sim P(x))$
- 3. Prove by direct proof
- 3.1 Suppose x is an element of D s.t. Q(X) is false
- 3.2 Show that P(x) is false.
- 4. Therefore, original statement is true

# 4 Set Theory

**Defn.** Set: Unordered collection of objects

Order and duplicates don't matter

**Defn.** Membership of Set  $\in$ : If S is set,  $x \in S$  means x is an element of S

**Defn.** Cardinality of Set |S|: The number of elements in S

Common Sets:

 $\mathbb{N}$  - Natural Numbers,  $\{0, 1, 2\}$ 

 $\mathbb{Z}$  - Integers

Q - Rational

 $\mathbb{R}$  - Real

 $\mathbb C$  - Complex

 $\mathbb{Z}^{\pm}$  - Positive/Negative Integers

**Defn.** Subset  $A \subseteq B \Leftrightarrow$  Every element of A is also an element of B  $A \subseteq B \Leftrightarrow \forall x (x \in A \Rightarrow x \in B)$ 

**Defn.** Proper Subset  $A \subsetneq B \Leftrightarrow (A \subseteq B \land A \neq B)$ 

**Theorem 6.2.4.** An empty set is a subset of every set, i.e.  $\emptyset \subseteq A$  for all sets A

**Defn.** Cartesian Product  $A \times B = \{(a, b) : a \in A \land b \in B\}$ 

**Defn.** Set Equality  $A = B \Leftrightarrow A \subseteq B \land B \subseteq A$  $A = B \Leftrightarrow \forall x (x \in A \Leftrightarrow x \in B)$ 

**Defn.** Union:  $A \cup B = \{x \in U : x \in A \lor x \in B\}$ 

**Defn.** Intersection:  $A \cap B = \{x \in U : x \in A \land x \in B\}$ 

**Defn.** Difference:  $B \setminus A = \{x \in U : x \in B \land x \notin A\}$ 

**Defn.** Disjoint:  $A \cap B = \emptyset$ 

**Theorem 4.4.1.** Quotient-Remainder  $n \in \mathbb{Z}, d \in \mathbb{Z}^+$ 

there exists unique integers q and r such that n = dq + r and  $0 \le r < d$ 

**Defn.** Power Set: The set of all subsets of A, has  $2^n$  elements.

**Theorem 6.3.1.** Suppose A is a finite set with n elements, then P(A) has  $2^n$  elements.  $|P(A)| = 2^{|n|}$ 

**Defn.** Cartesian Product of  $A_n = A_1 \times A_2 \times ... \times A_n = \{(a_1, a_2, ... a_n) : a_1 \in A_1 \wedge a_2 \in A_2...$ 

Theorem 6.2.1. Subset Relations

- 1. Inclusion of Intersection:  $A \cup B \subseteq A, A \cup B \subseteq B$
- 2. Inclusion in Union  $A \subseteq A \cup B, B \subseteq A \cup B$
- 3. Transitive Property of Substs:  $A \subseteq B \land B \subseteq C \Rightarrow A \subseteq C$

# 5 Relations

**Defn.** Relation from A to B is a subset of  $A \times B$ 

Given an ordered pair  $(x,y) \in A \times B$ , x is related to y by R is written  $xRy \Leftrightarrow (x,y) \in R$ 

Defn. Domain, Co-domain, Range

Let A and B be sets and R be a relation from A to B

- 1. Domain of R: is set  $\{a \in A : aRb \text{ for some } b \in B\}$
- 2. Codomain of R: Set B
- 3. Range of R: is set  $\{b \in B : aRb \text{ for some } a \in A\}$

**Defn.** Inverse Relation

Let R be a relation from A to B,  $R^{-1} = \{(y, x) \in B \times A : (x, y) \in R\}$  $\forall x \in A, \forall y \in B((y, x) \in R^{-1} \Leftrightarrow (x, y) \in R)$ 

**Defn.** Relation on a Set A is a relation from A to A.

**Defn.** Composition of Relations

A, B and C be sets.  $R \subseteq A \times B$  be a relation.  $S \subset B \times C$  be relation. Composition of R with S, denoted  $S \circ R$  is relation from A to C such that:

 $\forall x \in A, \forall z \in C(xS \circ Rz \Leftrightarrow (\exists y \in B(xRy \land ySz)))$ 

**Proposition.** Composition is Associative A, B, C, D be sets.  $R \subseteq A \times B, S \subseteq B \times C, T \subseteq C \times D$   $T \circ (S \circ R) = T \circ S \circ R$ 

**Proposition.** Inverse of Composition A,B,C be sets.  $R\subseteq A\times B, S\subseteq B\times C$   $(S\circ R)^{-1}=R^{-1}\circ S^{-1}$ 

### Defn. Reflexivity, Symmetry, Transitivity

- 1. Reflexivity:  $\forall x \in A(xRx)$
- 2. Symmetry:  $\forall x, y \in A(xRy \Rightarrow yRx)$
- 3. Transitivity: $\forall x, y, z \in A(xRy \land yRz \Rightarrow xRz)$

Refer to proof 6

**Defn.** Transitive Closure

Transitive closure of R is relation  $R^t$  on A that satisfies

- 1.  $R^t$  is transitive
- 2.  $R \subseteq R^t$
- 3. If S is any other transitive relation that contains R, then  $R^t \subseteq S$

Defn. Partition

P is partition of set A if

- 1. P is a set of which all elements are non empty subsets of A,  $\emptyset \neq S \subseteq A$  for all  $S \in P$
- 2. Every element of A is in exactly on element of P,

 $\forall x \in A \ \exists S \in P(x \in S) \ \text{and}$ 

 $\forall x \in A \ \exists S_1, S_2 \in P(x \in S_1 \land x \in S_2 \Rightarrow S_1 = S_2)$ 

OR  $\forall x \in A \ \exists ! S \in P(x \in S)$ 

Elements of a partition are called components

**Defn.** Relation Induced by a partition

Given partition P of A, the relation R induced by partition:

 $\forall x, y \in A, xRy \Rightarrow \exists$  a component of S of P s.t.  $x, y \in S$ 

**Theorem 8.3.1** (Relation Induced by a Partition). Let A be a set with a partition and let R be a relation induced by the partition. Then R is reflexive, symmetric and transitive

**Defn** (Equivalence Relation). A be set and R be relation. R is equivalence relation iff R is reflexive, symmetric and transitive

#### **Defn.** Equivalence Class

Suppose A is set and  $\sim$  is equivalence relation on A. For each  $A \in A$ , equivalence class of a, denoted [a] and called class of a is set of all elements  $x \in A$  s.t.  $a \sim x$ 

 $[a]_{\sim} = \{ x \in A : a \sim x \}$ 

#### **Theorem 8.3.4.** The partition induced by an Equivalence Relation

If A is a set and R is an equivalence relation on A, then distinct equivalence classes of R form a partition of A; that is, the union of the equivalence classes is all of A, and the intersection of any 2 discrinct classes is empty.

#### Defn. Congruence

Let  $a, b \in \mathbb{Z}$  and  $n \in \mathbb{Z}^+$ . Then a is congruent to b modulo n iff a - b = nk, for some  $k \in \mathbb{Z}$ . In other words, n | (a - b). We write  $a \equiv b \pmod{n}$ 

#### **Defn.** Set of equivalence classes

Let A be set and  $\sim$  be an equivalence relation on A. Denote by  $A/\sim$ , the set of all equivalence classes with respect to  $\sim$ , i.e.

 $A/\sim = \{[x]_\sim : x \in A\}$ 

**Theorem Equivalence Classes.** form a partition Let  $\sim$  be an equiv. relation on A. Then  $A/\sim$  is a partition of A.

**Defn** (Antisymmetry). R is antisymmetric iff  $\forall x, y \in A(xRy \land yRx \Rightarrow x = y)$  (DOES NOT IMPLY NOT SYMMETRIC)

**Defn** (Partial Order Relation). R is Partial Order iff R is reflexive, antisymmetric and transitive.

**Defn.** Partially Ordered Set Set A is called poset with respect to partial order relation R on A, denoted by (A, R) (Proof 7)

**Defn.**  $x \leq y$  is used as a general partial order relation notation

**Defn** (Hasse Diagram). Let  $\leq$  be a partial order on set A. Hasse diagram satisfies the following condition for all distinct  $x, y, m \in A$ 

If  $x \leq y$  and no  $m \in A$  is s.t.  $x \leq m \leq y$ , then x is placed below y with a line joining them, else no line joins x and y.

**Defn** (Comparability).  $a, b \in A$  are comparable iff  $a \leq b$  or  $b \leq a$ . Otherwise, they are **noncomparable** 

**Defn** (Maximal, Minimal, Largest Smallest). Set A be partially ordered w.r.t. a relation  $\leq$  and  $c \in A$ 

- 1. c is maximal element of A iff  $\forall x \in A$ , either  $x \leq c$  or x and c are non-comparable. OR  $\forall x : nA(c \leq x \Rightarrow c = x)$
- 2. c is minimal element of A iff  $\forall x \in A$ , either  $c \preceq x$  or x and c are non-comparable. OR  $\forall x in A (x \preceq c \Rightarrow c = x)$
- 3. c is largest element of A iff  $\forall x \in A(x \leq c)$
- 4. c is smallest element of A iff  $\forall x \in A(c \leq x)$

## **Proposition.** A smallest element is minimal

Consider a partial order  $\leq$  on set A. Any smallest element is minimal.

- 1. Let c be smallest elemnt
- 2. Take any  $x \in A$  s.t.  $x \leq c$
- 3. By smallestness, we know  $c \leq x$  too.
- 4. So c = x by antisymmetry

**Defn** (Total Order Relations). All elements of the set are comparable

R is total order iff R is a partial order and  $\forall x, y \in A(xRy \vee yRx)$ 

**Defn** (Linearization of a partial order). Let  $\leq$  be a partial order on set A. A linearization of  $\leq$  is a total order  $\leq$  \* on

 $A \text{ s.t. } \forall x, y \in A(x \leq y \Rightarrow x \leq y)$ 

**Defn** (Kahn's Algorithm). Input: A finite set A and partial order  $\leq$  on A

- 1. Set  $A_0 := A$  and i := 0
- 2. Repeat until  $A_i = \emptyset$ 
  - 2.1. Find minimal element  $c_i$  of  $A_i$  wrt  $\leq$
  - 2.2. Set  $A_{i+1} = A_i \setminus c_i$
  - 2.3. Set i = i + 1

Output: A linearization  $\leq *$  of  $\leq$  defined by setting, for all indicies i, j  $c_i \leq *$   $c_j \Leftrightarrow i \leq j$ 

**Defn** (Well ordered set). Let  $\leq$  be a total order on set A. A is well ordered iff every nonempty subset of A contains a smallest element. OR

 $\forall S \in P(A), S \neq \emptyset \Rightarrow (\exists x \in S \forall y \in S(x \leq y)) \text{ E.g. } (\mathbb{N}, \leq) \text{ is well ordered but } (\mathbb{Z}, \leq) \text{ is not as there is no smallest integer (Theorem 4.6.1)}$ 

## 6 Proofs

Proof L1S28. Prove that the product of two consecutive odd numbers is always odd.

- 1. Let a and b be two consecutive odd numbers
- 1.1. Without loss of generality, assume that a < b, hence b = a + 2
- 1.2. Now, a = 2k + 1 (by defin of odd numbers)
- 1.3. Similarly, b = a + 2 = 2k + 3
- 1.4. Therefore,  $ab = (2k+1)(2k+3) = (4k^2+6k) + (2k+3) = 4k^2+8k+3 = 2(2k^2+4k+1)+1$  (by Basic Algebra)
- 1.5. Let  $m = (2k^2 + 4k + 1)$  which is an integer (by closure of integers under  $\times$  and +)
- 1.6. Then ab = 2m + 1 which is odd (by defin of odd numbers)
- 2. Therefore, the product of two consecutive odd numbers is always odd.

Proof L4S16. Sum of 2 even  $\mathbb{Z}$  is even

- 1. Let m and n be two particular but arbitrarily chosen even intergers
- 1.1. Then m=2r and n=2s for some  $\mathbb{Z}$  r and s (by defining of even number)
- 1.2. m + n = 2r + 2s = 2(r + s) (by basic algebra)
- 1.3. 2(r+s) is an integer (closure of int under  $\times$  and +) and an even number (by defin of even number)
- 1.4. Hence m + n is an even number
- 2. Therefore sum of any two even integers is even

Proof T 4.6.1. There is no greatest integer (Contradiction)

- 1. Suppose not, i.e. there is a greatest intger
- 1.1. Lets call this greatest integer g, and  $g \geq n$  for all integers n
- 1.2. Let G = g + 1
- 1.3. Now, G is an integer (closure of integers under +) and G > g
- 1.4. Hence, g is not the greatest integer, contradicting 1.1
- 2. Hence, the supposition that there is a greatest integer is false.
- 3. Therefore there is no greatest integer

Proof L5S19. L5S19 Two sets are equal

- 1. Let sets X and Y be given. To prove X = Y
- 2. ( $\subseteq$ ) Prove  $X \subseteq Y$
- 3.  $(\supseteq)$  Prove  $X \supseteq Y$
- 4. From (2) and (3), we can conclude that X = Y

# Proof L5S22. L5S22 $\{x \in Z : x^2 = 1\} = \{1, -1\}$ 1.1. Take any $z \in \{x \in \mathbb{Z} : x^2 = 1\}$ 1.2. Then $z \in \mathbb{Z}$ and $z^2 = 1$ 1.3. So, $z^2 - 1 = (z - 1)(z + 1) = 0$ (by basic algebra) 1.4. $\therefore z - 1 = 0 \text{ or } z + 1 = 0$ 1.5. $\therefore z = 1 \text{ or } z = -1$ 1.6. So, $z \in \{1, -1\}$ $2. \leftarrow$ 2.1. Take any $z \in \{1, -1\}$ 2.2. Then z = 1 or z = -12.3. In either case, we have $z \in \mathbb{Z}$ and $z^2 = 1$ 2.4. So, $z \in \{x \in \mathbb{Z} : x^2 = 1\}$ 3. Therefore, $\{x \in Z : x^2 = 1\} = \{1, -1\}$ (from (1) and (2)) Proof L6S27. $\forall x,y \in \mathbb{Z}(xRy \Leftrightarrow 3|(x-y))$ is reflexive, symmetric, transitive 1. Proof of Reflexivity 1.1. Let a be an arbitrarily chosen integer. 1.2. Now a - a = 01.3. 3|0(since $0 = 3 \cdot 0$ ), hence 3|(a - a)1.4. Therefore aRa (by defin of R) 2. Proof of Symmetry 2.1. Let a, b be arbitrarily chosen integers 2.2. Then 3|(a-b) (by defin of R), hence a-b=3k for some integer k (by defin of divisibility) 2.3. Multiplying both sides by -1 gives b - a = 3(-k)2.4. Since -k is an integer, 3|(b-a) (by defin of divisibility) 2.5. Therefore, $aRb \Rightarrow bRa$ (by defn of R) 3. Proof of Transitivity 3.1. Let a, b, c be arbitrarily chosen integers 3.2. Then, 3|(a-b) and 3|(b-c) (by defin of R), hence a-b=3r and b-c=3s (by defin of divisibility) 3.3. Adding both equations gives a - c = 3r + 3s3.4. Since r + s is an integer, 3|(a - c) (by defin of divisibility) 3.5. Therefore $aRb \wedge bRc \Rightarrow aRc$ (by defin of R) Lemma Equivalence Class L6S47. Let $\sim$ be an equivalence relation on A. The following are equivalent for all $x, y \in A$ (i) $x \sim y$ , (ii) [x] = [y], (iii) $[x] \cap [y] \neq \emptyset$ 1. $x \sim y \Rightarrow [x] = [y]$ 1.1. Suppose $x \sim y$ 1.2. Then $y \sim x$ (by symmetry) 1.3. For every $z \in [x]$ 1.3.1. $x \sim z$ (by defin of x) 1.3.2. $\therefore y \sim z$ (by transitivity of $y \sim x$ ) 1.3.3. $\therefore z \in [y]$ (by defin of [y]) 1.4. This shows $[x] \subseteq [y]$ 1.5. Switching roles of x and y, we can also see that $[y] \subseteq [x]$ 1.6. Therefore, [x] = [y] $2. \quad [x] = [y] \Rightarrow [x] \cap [y] \neq \emptyset$ 2.1. Suppose [x] = [y]2.2. Then $[x] \cap [y] = [x]$ (by idempotent law for $\cap$ ) 2.3. However, we know $x \sim x$ (by reflexivity of $\sim$ ) 2.4. This shows $x \in [x] = [x] \cap [y]$ (by defin of [x] and (2.2)) 2.5. Therefore $[x] \cap [y] \neq \emptyset$ 3. $[x] \cap [y] \neq \emptyset \Rightarrow x \sim y$ 3.1. Suppose $[x] \cap [y] \neq \emptyset$ 3.2. Take $z \in [x] \cap [y]$ 3.3. Then $z \in [x]$ and $z \in [y]$ (by defin of $\cap$ )

3.4. Then  $x \sim z$  and  $y \sim z$  (by defin of [x] and [y])

- 3.5.  $y \sim z$  implies  $z \sim y$  (by defin of symmetry)
- 3.6. Therefore,  $x \sim y$  (by transitivity)

Proposition L6S54. Congruence-mod n is an equivalence relation on  $\mathbb{Z}$  for every  $n \in \mathbb{Z}^+$ 

- 1. (Reflexivity) For all  $a \in \mathbb{Z}$
- 1.1.  $a a = 0 = n \times 0$
- 1.2. So  $a \equiv a \pmod{n}$  (by definition of congruence)
- 2. (Symmetry)
- 2.1. Let  $a, b \in \mathbb{Z}$  s.t.  $a \equiv a \pmod{n}$
- 2.2. Then there is a  $k \in \mathbb{Z}$  s.t. a b = nk
- 2.3. Then b a = -(a b) = -nk = n(-k)
- 2.4.  $-k \in \mathbb{Z}$  (by closure of integers under  $\times$ ), so  $b \equiv a \pmod{n}$  (by defin of congruence)
- 3. (Transitivity)
- 3.1. Let  $a, b, c \in \mathbb{Z}$  s.t.  $a \equiv a \pmod{n}$  and  $b \equiv c \pmod{n}$
- 3.2. Then there is a  $k, l \in \mathbb{Z}$  s.t. a-b=nk and b-c=nl
- 3.3. Then a c = (a b) + (b c) = nk + nl = n(k + 1)
- 3.4.  $k+l \in \mathbb{Z}$  (by closure of integers under +), so  $a \equiv c \pmod{n}$  (by defin of congruence)

Proof L6S69.  $\forall a, b \in \mathbb{Z}^+, \forall a | b \Leftrightarrow b = ka$  for some integer k. Prove | is a partial order relation on A

- 1. | is reflexive: Suppose  $a \in A$ . Then  $a = 1\dot{a}$ , so a|a (by defining of divisibility)
- 2. | is antisymmetric
- 2.1. Suppose  $a, b \in \mathbb{Z}^+$  such that aRb and bRa
- 2.2. Then b = ra and a = sb for some integers r and s (by defin of divides). It follows that b = ra = r(sb)
- 2.3. Dividing both sides by b gives 1 = rs
- 2.4. Only product of two positive integers that equals 1 is 11.
- 2.5. Thus r = s = 1, and so  $a = sb = 1\dot{b} = b$
- $2.6. \quad \text{Therefore,} \mid \text{is antisymmetric}$

OR

- 2.1. Suppose  $a, b \in \mathbb{Z}^+$  such that a|b and b|a
- 2.2. then  $a \leq b$  and  $b \leq a$  (by theorem 4.3.1)
- 2.3. So a = b
- 3. | is transitive: Show that  $\forall a, b, c \in A, a | b \land b | c \Rightarrow a | c$ ) (theorem 4.3.3)

*Proof T01Q9*. The product of any two odd integers is an odd integer

- 1. Take any 2 odd numbers a and b
- 2. Then a = 2k + 1 and b = 2p + 1 for  $k, p \in \mathbb{Z}$  (by defin of odd number)
- 3. Then  $a \cdot b = (2k+1)(2p+1) = (4kp+2k) + (2p+1) = 2(2kp+p+k) + 1$  (by defin of odd number)
- 4. Let q = 2kp + p + k which is an integer (by closure of int under + and ×
- 5. Then nm = 2q + 1 which is odd (by defin of odd numbers)

*Proof T01Q10.* Let n be an integer. Then  $n^2$  is odd iff n is odd

- 1. Proof By Contraposition, that is "if n is even,  $n^2$  is even  $(\Rightarrow)$
- 1.1. Suppose n is even.
- 1.2. Then  $\exists k \in \mathbb{Z} \text{ s.t. } n = 2k \text{ (by defn of even integers)}$
- 1.3.  $n^2 = (2k)^2 = 4k^2 = 2(2k^2)$
- 1.4. Hence,  $n^2 = 2p$ , where  $p = 2k^2 \in \mathbb{Z}$  (by closure of integers under  $\times$ )
- 1.5. Therefore,  $n^2$  is even and this proves that if  $n^2$  is odd, n is odd.
- 2. If n is odd, then  $n \times n = n^2$  is odd (T01Q9)
- 3. Therefore  $n^2$  is odd if and only if n is odd.

## $Proof\ T02Q3$ . Rational numbers are closed under addition

- 1. Let r and s be rational numbers
- 2.  $\exists a, b, c, d \in \mathbb{Z}$  s.t.  $r = \frac{a}{b}, s = \frac{c}{d}$  and  $b \neq 0, d \neq 0$  (by defin of rational numbers)
- 3. Hence  $r + s = \frac{a}{b} + \frac{c}{d} = \frac{ad + bc}{bd}$  (by basic algebra)
- 4.  $ad + bd \in Z$  and  $bd \in Z$  (closure of integers under + and  $\times$ )
- 5.  $bd \neq 0$  since  $b \neq 0, d \neq 0$
- 6. Hence r + s is rational, therefore rational numbers are closed under addition

*Proof T02Q10.* if n is a product of 2 positive integers a and b, then  $a \leq n^{1/2}$  or  $b \leq n^{1/2}$ 

- 1. Proof by contraposition, that is if  $a > n^{1/2}$  and  $b > n^{1/2}$ , then n is not a product of a and b
- 2. Suppose  $a > n^{1/2}$  and  $b > n^{1/2}$ , then  $ab > n^{1/2} \cdot n^{1/2} = n$  (by Appendix A T27)
- 3. Since  $ab \neq n$ , the contrapositive statement is true

or by contradiction

- 1. Proof by contradiction, that is n = ab and  $a > n^{1/2}$  and  $b > n^{1/2}$
- 2. Since  $a > n^{1/2}$  and  $b > n^{1/2}$ , then  $ab > n^{1/2} \cdot n^{1/2} = n$  (by Appendix A T27)
- 3. This contradicts n = ab. Therefore original statement is true

*Proof T03Q04.* Let  $A = \{2n + 1 : n \in \mathbb{Z}\}$  and  $B = \{2n - 5 : n \in \mathbb{Z}\}$ . Is A = B?

- 1. ⊆
  - 1.1. Let  $a \in A$ , and  $a = 2n + 1, n \in \mathbb{Z}$
  - 1.2. Then a = 2n + 1 = 2(n+3) 5
  - 1.3.  $n+3 \in \mathbb{Z}$  (by closure of int under +)
  - 1.4. Therefore,  $a \in B$  (by defn of B)
- 2. ⊒
  - 2.1. Let  $b \in A$ , and  $b = 2n 5, n \in \mathbb{Z}$
  - 2.2. Then b = 2n 5 = 2(n 3) + 1
  - 2.3.  $n-3 \in \mathbb{Z}$  (by closure of int under -)
  - 2.4. Therefore,  $b \in A$  (by defin of B)
- 3. Therefore, A = B

*Proof T03Q05.* Prove  $\forall A, B, C, A \cap (B \setminus C) = (A \cap B) \setminus C$ 

- 1.  $A \cap (B \setminus C) = \{x : x \in A \land x \in (B \setminus C)\}$  (by defin of  $\cap$ )
- 2.  $= \{x : x \in A \land (x \in B \land x \notin C)\}$  (by defin of \)
- 3.  $= \{x : x \in (A \land x \in B) \land x \notin C\}$  (by associativity of  $\land$ )
- 4.  $= \{x : x \in (A \cap B) \land x \notin C\}$  (by defin of  $\cap$ )
- 5. =  $\{x : x \in (A \cap B) \setminus C \text{ (by defn of } \setminus \}$

*Proof T03Q05.* Prove  $\forall A, B, C, A \cap (B \setminus C) = (A \cap B) \setminus C$ 

- 1.  $A \cap (B \setminus C) = \{x : x \in A \land x \in (B \setminus C)\}$  (by defin of  $\cap$ )
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- 4.  $= \{x : x \in (A \cap B) \land x \notin C\}$  (by defin of  $\cap$ )
- 5.  $= \{x : x \in (A \cap B) \setminus C \text{ (by defn of } \setminus \}$

*ProofT03Q8.* Let A and B be set. Show that  $A \subseteq B$  if and only if  $A \cup B = B$  To show  $A \cup B = B$ , we need to show  $A \cup B \subseteq B$  and  $B \subseteq A \cup B$ 

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1.1. Suppose A \subseteq B
    1.2.
          (Show A \cup B \subseteq B)
           1.2.1. Let z \in A \cup B
           1.2.2. Then z \in A or z \in B (by defin of \cup)
          1.2.3. Case 1: Suppose z \in A, then Z \in B as A \subseteq B line (1.1)
           1.2.4. Case 2: Suppose z \in B, then z \in B. We have z \in B in either case
    1.3. (Show A \cup B \supset B)
           1.3.1. Let z \in B
           1.3.2. Then z \in A or z \in B (by generalization)
           1.3.3. So z \in A \cup B (by defin of \cup)
          Therefore A \cup B = B
    1.4.
    \Leftarrow
    2.1. Suppose A \cup B = B
    2.2. Let z \in A
           2.2.1. Then z \in A or z \in B (by generalization)
           2.2.2. So z \in A \cup B (by defin of \cup)
           2.2.3. So z \in B since A \cup B = B (2.1)
          Therefore A \subseteq B
   Therefore, A \subseteq B if and only iff A \cup B = B
    Proof T04Q05. Relation S = \{(m, n) \in \mathbb{Z}^2 : m^3 + n^3 \text{ is even}\}, Proof S \circ S = S
1. (\subseteq) Suppose (x, z) \in S \circ S
    1.1. Then (x,y) \in S and (y,z) \in S for some y \in Z (defin of composition of relations)
    1.2. So x^3 + y^3 is even and y^3 + z^3 is even
    1.3. This implies that x^3 + 2y^3 + z^3 is even
    1.4. This implies that x^3 + z^3 is even as 2y^3 is even
    1.5. Therefore, (x, z) \in S (by defin of S)
    (\supseteq) Suppose (x,z) \in S
    2.1. Then x^3 + z^3 is even (by defin of S)
          Case 1: x^3 is odd.
           2.2.1. Then z^3 is also odd.
           2.2.2. This implies that x^3 + 1^3 is even and 1^3 + z^3 is even
           2.2.3. Thus, (x, 1) \in S and (1, z) \in S (by defin of S)
           2.2.4. So, (x, z) \in S \circ S
    2.3. Case 2: x^3 is even.
          2.3.1. Then z^3 is also even.
          2.3.2. This implies that x^3 + 0^3 is even and 0^3 + z^3 is even
           2.3.3. Thus, (x,0) \in S and (0,z) \in S (by defin of S)
           2.3.4. So, (x, z) \in S \circ S
    2.4. In all cases, (x, z) \in S \circ S
    OR
   (\supseteq) Suppose (x, z) \in S
    3.1. Note that (x, x) \in S as x^3 + x^3 is even
    3.2. Since (x,x) \in S and (x,z) \in S, we have (x,z) \in S \circ S (by defin of composition of relations)
    Proof. R is asymmetric if and only if R is antisymmetric and irreflexive.
    1.
        \Longrightarrow
        1.1. R is irreflexive (R is irreflexive \implies R is antisymmetric and irreflexive)
               1.1.1. Let x \in A s.t. xRx \implies x \not Rx (R is Asymmetric)
               1.1.2. Since x \not R x, R is irreflexive (by defining of irreflexive)
         1.1. R is antisymmetric (Tutorial Qn 6c)
        \Leftarrow (R is antisymmetric and irreflexive \Longrightarrow asymmetry)
        2.1. Let x, y \in A, s.t. xRy is antisymmetric and irreflexive
        2.2. There is 2 cases to consider, x = y and x \neq y
        2.3. x = y
               2.3.1. xRx is not valid as it contradicts irreflexive, \forall x \in A(x Rx)
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1.

2.3.2. Therefore,  $xRx \implies x Rx$ 

$$\begin{array}{ccc} 2.4. & x \neq y \\ & 2.4.1. & xRy \wedge yRx \implies x = y \end{array}$$

# 7 Tables

Commutative	$p \wedge q \equiv q \wedge p$	$p \lor q \equiv q \lor p$
Associative	$p \wedge q \wedge r \equiv (p \wedge q) \wedge r$	
Distributive	$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	$p \lor (q \land r) \equiv (p \lor q) \land (p \lor r)$
Identity	$p \wedge \text{true} \equiv p$	$p \lor (q \lor \lor) = (p \lor q) \lor (p \lor \lor)$ $p \lor \text{false} \equiv p$
Negation	$p \lor crue = p$ $p \lor \sim p \equiv \text{true}$	$p \land \text{ranke} = p$ $p \land \sim p \equiv \text{false}$
_		$p \wedge p \equiv \text{raise}$
Double Negative	$\sim (\sim p) \equiv p$	
Idempotent	$p \lor p \equiv p$	$p \wedge p \equiv p$
Universal bound	$p \vee \text{true} \equiv \text{true}$	$p \land \text{false} \equiv \text{false}$
de Morgan's	$\sim (p \land q) \equiv \sim p \lor \sim q$	$\sim (p \lor q) \equiv \sim p \land \sim q$
Absorption	$p \lor (p \land q) \equiv p$	$p \land (p \lor q) \equiv p$
Implication	$p \Rightarrow q \equiv \sim p \lor q$	
$\sim$ (Implication)	$\sim (p \Rightarrow q) \equiv p \land \sim q$	
Modus Ponens	$p \implies q, p$	q
Modus Tollens	$p \implies q, \sim q$	$\sim p$
Generalization	p $p$	$p \lor q$
Specialization	$p \land q$	
		p
Conjunction	p,q	$p \wedge q$
Elimination	$p \lor q, \sim q$	p
Transitivity	$p \implies q, q \implies r$	$p \implies r$
Division into cases	$p \wedge q, p \implies r, q \implies r$	r
Contradiction	$\sim p \implies \text{false}$	p
Commutative	$A \cup B = B \cup A$	
Associative	$(A \cup B) \cup C = A \cup (B \cup C)$	
Distributive	$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$	$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$
Identity	$A \cup \emptyset = A$	$A \cap U = A$
Complement	$A \cup \bar{A} = U$	$A \cap \bar{A} = \emptyset$
_	$ar{ar{A}} = A$	$21 + 21 = \psi$
Double Complement		4 0 4 4
Idempotent	$A \cup A = A$	$A \cap A = A$
Universal Bound	$A \cup U = U$	$A \cap \emptyset = \emptyset$
De Morgan's	$\overline{A \cup B} = \overline{A} \cap \overline{B}$	$\overline{A \cap B} = \bar{A} \cup \bar{B}$
Absorption	$A \cup (A \cap B) = A$	$A \cap (A \cup B) = A$
Complements of U and $\emptyset$	$\bar{U} = \emptyset$	$ar{\emptyset} = U$
Set Difference	$A \setminus B = A \cap \bar{B}$	
	,	
F1 Commutative	a+b=b+a	ab = ba
F2 Associative	(a+b)+c=a+(b+c)	(ab)c = a(bc)
F3 Distributive	a(b+c) = ab + ac	(b+c)a = ba + ca
F4 Identity	0 + a = a + 0 = a	$1 \cdot a = a \cdot 1 = a$
· ·		$1 \cdot a = a \cdot 1 = a$
F5 Additive inverses	a + (-a) = (-a) + a = 0	- / 0
F6 Reciprocals	$a \cdot \frac{1}{a} = \frac{1}{a} \cdot a = 1$	$a \neq 0$
T1 Concellation Add	a+b=a+a	b=c
T1 Cancellation Add	a+b=a+c	b = c
T2 Possibility of Sub	There is one $x, a + x = b$	x = b - a
T3	b - a = b + (-a)	
T4	-(-a) = a	
T5	a(b-c) = ab - ac	
Т6	$0 \cdot a = a \cdot 0 = 0$	
T7 Cancellation Mul	ab = ac	$b = c, a \neq 0$
T8 Possibility of Div	$a \neq 0, ax = b$	$x = \frac{b}{a}$
T9	$a \neq 0, \frac{b}{a} = b \cdot a^{-1}$	a
T10	$a \neq 0, \frac{a}{a} = 0 \cdot a$ $a \neq 0, (a^{-1})^{-1} = a$	
	$a \neq 0, (a  ) = a$ $ab = 0 \Rightarrow a = 0 \lor b = 0$	
T11 Zero Product		a - a - a
T12 Mul with -ve	(-a)b = a(-b)(ab)	$-\frac{a}{b} = \frac{-a}{b} = \frac{a}{-b}$
T13 Equiv Frac	$\frac{a}{b} = \frac{ac}{bc}$	$b \neq 0, c \neq 0$
L 001 4 A 1 1 D	$\underline{a} \perp \underline{c} = \underline{ad+bc}$	$b \neq 0, d \neq 0$
T14 Add Frac	b + d - bd	1 , 0, 11 , 0
T14 Add Frac T15 Mul Frac	$ \begin{array}{ccc} b & d & bd \\ \frac{a}{h} \cdot \frac{c}{d} & \frac{ac}{hd} \end{array} $	$b \neq 0, d \neq 0$
	$\frac{a}{b} = \frac{ac}{bc}$ $\frac{a}{b} + \frac{c}{d} = \frac{ad + bc}{bd}$ $\frac{a}{b} \cdot \frac{c}{d} = \frac{ac}{bd}$ $\frac{\frac{a}{b}}{\frac{c}{d}} = \frac{ac}{bd}$	1 1 1

Ord1	$\forall a,b \in \mathbb{R}^+$	a+b>0, ab>0
Ord2	$\forall a, b \in \mathbb{R}_{\neq 0}$	a is positive or negative and not both
Ord3	0 is not positive	
a < b	means $b + (-a)$ is positive	
$a \leq b$	means $a < b$ or $a = b$	
a < 0	means a is negative	
T17 Trichotomy Law	$a < b \lor b > a \lor a = b$	
T18 Transitive Law	a < b and $b < c$	a < c
T19	a < b	a + c < b + c
T20	a < b  and  c > 0	ac < bc
T21	$a \neq 0$	$a^2 > 0$
T22	1 > 0	
T23	a < b and $c < 0$	ac > bc
T24	a < b	-a > -b
T25	ab > 0	a and b are both positive or negative
T26	a < c and $b < d$	a+b < c+d
T30	0 < a < c  and  < 0 < b < d	0 < ab < cd